# Multiforked Iterated Even-Mansour and a Note on the Tightness of IEM Proofs

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Selected Areas in Cryptography 2025, Toronto, Canada, 13. Aug 2025





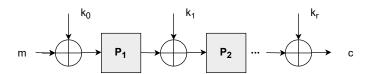
## Iterated Even-Mansour (IEM)

### many ciphers (e.g. AES):

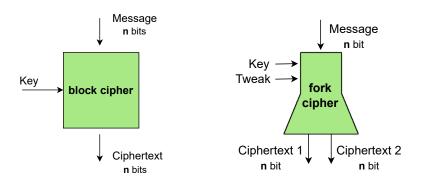
- repeated round function
- key expanded into round keys

#### IEM:

- ightharpoonup public permutations  $P_1, \ldots, P_r$
- $ightharpoonup k_0, \ldots, k_r$  uniformly random (idealized key schedule)

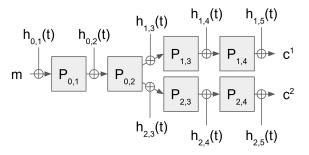


## Forkcipher



forkcipher applications: encryption [ABPV21], AEAD [ALP $^+$ 19], PRG [AW23], KDF [BDA $^+$ 24], ...

## Forked IEM (Our work)

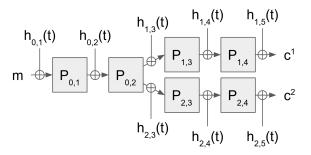


Forked IEM (4 rounds, 2 branches)

#### **Variants**

- ▶ no tweak:  $h_i(\cdot)$  returns round key  $k_i$
- ▶ idealized tweakey schedule:  $h_i(\cdot)$  = random function

## Forked IEM (Existing Variant)



Forked IEM (4 rounds, 2 branches)

#### **Variants**

- AXU tweakey schedule [KLL20]:  $h_i(\cdot)$  based on AXU hash existing proof [KLL20]: only 2 rounds
  - $\Rightarrow$  our proof: arbitrary rounds and branches

# Security of IEM Variants

Tweakey schedule	IEM/TEM	Forked IEM
no tweaks	$2^{r  n/(r+1)}$ [HT16]	$2^{r n/(r+1)}$ [our work]
idealized	_	$2^{r n/(r+1)}$ [our work]
AXU (2 rounds)	$2^{r  n/(r+1)}$ [CLS15]	$2^{r  n/(r+1)}$ [KLL20]
AXU (unrestricted)	$2^{r  n/(r+2)}$ [CLS15]	$2^{r n/(r+2)}$ [our work]

Security (in queries). *r* rounds construction.

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#### More than 2 branches?

**b** branches (AXU schedule, r rounds):  $\frac{1}{b^2}2^{r n/(r+2)}$  queries

## **Proof Approach**

- no tweaks: Expectation method [HT16]
  - represent attacker knowledge as graph & simplify graph
  - ▶ at the core: bound difference between 1 forked and 2 non-forked instances

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  - ▶ multi-user no tweak ≈ single-user ideal tweakey schedule

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- AXU tweak: Coupling [CLS15]
  - extending existing proof for non-forked to arbitrary many branches

## Tightness of IEM Proofs

- ▶ tightness: security proof + attack (practical efficiency!)
- ▶ unproven attack [BKL<sup>+</sup>12] used to argue tightness (directly or indirectly) [CLS15, BKL<sup>+</sup>12, LPS12, Ste12, CS14]

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We show: attack by Gaži [Gaž13] applies to IEM

⇒ tightness results remain

#### Conclusion

#### Main result: Forked IEM security

- arbitrary number of rounds
- 3 variants for tweakey schedule (no tweak / idealized / AXU)
- ightharpoonup security of forked IEM pprox non-forked IEM (with similar tweakey schedule)
- generalization to arbitrary number of branches for AXU variant

#### Note on tightness

▶ instantiation of Gaži [Gaž13] attack

## Thank you!

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